

## Prove that the Steinhaus method is always fair for $N = 3$

### Proof 3:

Easy case: A, B, C all make a bid list.  $N = 3$ . Each player is then given a piece from their bid list (which is composed of preferences that would give a fair share).

Hard case: B and C both only want  $S_1$ . We give A  $S_2$  or  $S_3$ , say  $S_3$ , and have A and B do I cut you choose on the recombined  $S_1 + S_2$ . A and B both value  $S_1 > 2/3$ , because  $S_2$  and  $S_3$  would not be fair to them. When they perform I cut you choose, each is guaranteed at least  $1/2$  of  $2/3$  which is  $1/3$ , a fair division.

should say B and C  
both value  $(S_1 + S_2) > 2/3$   
since  $S_3$  worth  $< 1/3$ ,  $(S_1 + S_2)$  worth  $> 2/3$

### Proof 5:

Easy case: Everyone gets a piece on their bid list.

Hard case: Let's say B and C both want  $S_1$  and only  $S_1$ . Let's say D wants  $S_2$  (or  $S_3$ ). D gets  $S_2$  - fair share. B and C believe  $S_1 > 1/3$   $S_2 < 1/3$   $S_3 < 1/3$ . Recombine  $S_1$  and  $S_3$  (or  $S_2$ ) and divide - "I cut, you choose". Both B and C will get half of piece that they believe is  $> 1/3$  - so it is fair to B and C too.

D lists all 3 on bidlist, so "wants" them all.  
We can choose to give D  $S_2$ .

B and C believe the recombined piece is worth  $> 2/3$

### Proof 7:

Easy case: Each player get piece of their bid list, so each gets at least  $1/3$ . Fair.

Hard case: - D divides 3 equal pieces, guaranteed to get at least  $1/3$ .

- B and D both only value  $S_1$
- This means they see  $S_1 > 2/3$ , while  $S_2$  &  $S_3$  both  $< 1/3$ .
- Each guaranteed  $1/2$  of  $2/3$ , which is  $1/3$ .
- Therefore, fair.

D gets exactly a fair share

Should say B and C

They only value  $S_1 > 1/3$ , not  $2/3$

Why? need to give a piece to D, describe recombined piece, and its worth, and then mention "I cut you choose", which is what guarantees  $1/2$  to each.